Security Vulnerabilities in LoRaWAN

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- Introduction
- Security features of LoRaWAN
- Attacks and attack mitigations
 - Replay attack
 - Eavesdropping
 - Bit-flipping attack
 - $\circ \quad \text{ACK spoofing} \quad$
 - LoRa class B attacks

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IoT Devices and Security

• Insecure defaults or insecure, remotely exploitable code

• Lack of accepted reference architecture

• Deployed in exposed and complex environment

LoRaWAN

• LoRaWAN: Long Range Wide Area Network

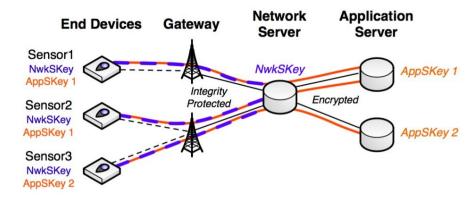
• MAC Layer protocol

• Widely adopted in IoT

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Channel Confidentiality

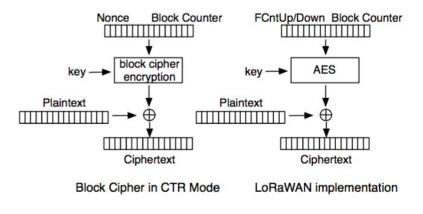
2 keys to ensure confidentiality:



Network Key (NwkSKey) \rightarrow between the IoT device and the network infrastructure

Application Key (AppSKey) \rightarrow between the IoT device and a third-party application provider in the backend

Channel Confidentiality - Message Encryption



LoRaWAN borrows idea from Block Cipher in CTR mode AppSKey encryption procedure:

- 1. Use **FCntUp** and **FCntDown** as Nonce;
- 2. The block cipher (AES) generate a pseudo-random permutation, which is used as a keystream;
- 3. Apply exclusive OR (XOR) to encrypt the plaintext.

Enrollment Protocol

The 2 keys, NwkSKey and AppSKey, needs to be 'enrolled' into the protocol.

By:

- 1. OTAA (Over-the-Air Activation)
- 2. ABP (Activation by Personalization)

Enrollment Protocol - OTAA

End device: send a *Join Request* (contains a 3-byte DevNonce) to network server;

Network server checks if 1) not accepted, then do nothing; 2) accepted, then sends a *Join Accept* to the end device.

Join Accept: contains a 3-byte AppNonce

Then both sides use this AppNonce to generate the NwkSKey and AppSKey

NwkSKey = AES_E (AppKey, 0x01 || AppNonce || NetID || DevNonce || pad) AppSKey = AES_E (AppKey, 0x02 || AppNonce || NetID || DevNonce || pad)

Enrollment Protocol - ABP

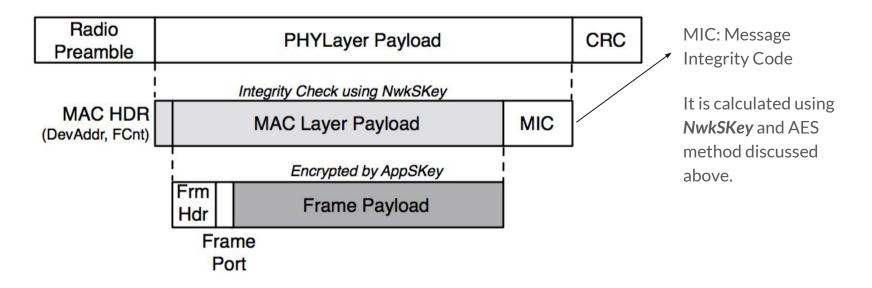
Different from OTAA: it skips the exchange of join messages;

DevAddr, NwkSKey, AppSKey are assigned to end devices, and stored in the server;

These unique parameters will be used across the sessions until updated in the device;

Encrypted messages are sent directly from an end device to server.

Integrity and Authenticity Validation



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Targets of Attacks

- 1. Adversary wants to *eavesdrop* and *decrypt* the content of a frame
- 2. It is possible that the content of a packet would be *modified outside of integrity check*
- 3. Messages would be *replayed*
- 4. A node would be tricked into believing that a message has been received successfully while it hasn't
- 5. Battery exhaustion attack

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Replay Attack for ABP-activated Nodes

Recall: ABP used *static* keys which are programmed into the device

Consequence: ABP-activated end device *reuses* the frame counter value from 0 with the *same* keys

- When does this happen?
- When counter overflows and is reset to 0.

Replay Attack for ABP-activated Nodes (Cont'd)

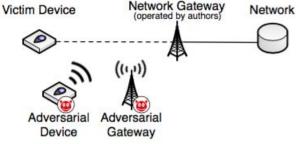
Steps for adversary to perform an attack:

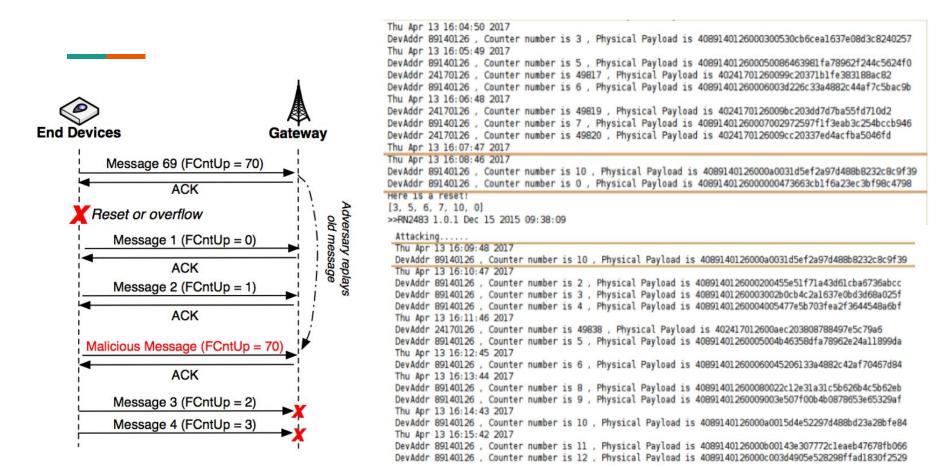
1. Monitor, store, keep in memory the uplink messages

2.Wait and prepare for the counter to be reset

3.Replay the malicious message when the time comes

*4.Replay this message again and again to block end device permanently (DoS attack)





_	time	counter	port	dev id	
	16:16:00	13	6	22	34 34 37 20 30 32 34 00
	1 6:15:25	12	61	22	34 39 36 20 30 32 34 00
	1 6:14:51	11	20	22	35 34 33 20 30 32 31 00
	1 6:08:49	10	49	22	34 38 30 20 30 32 31 00
	1 6:08:34	0	71	22	31 39 32 20 30 32 32 00
	▲ 16:07:59	10	49	22	34 38 30 20 30 32 31 00
	16:06:16	7	41	22	35 32 37 20 30 32 33 00
	▲ 16:05:42	6	61	22	36 38 37 20 30 32 34 00
	16:05:07	5	134	22	34 39 34 20 30 32 33 00
	1 6:03:59	3	83	22	34 34 38 20 30 32 32 00

Replay Attack Mitigation

Minimize the use of ABP, and use OTAA if possible

If insist on using ABP:

1.Adopt new keys periodically

2.Protect end devices physically, i.e. use non-volatile memory to preserve the counter value and avoid sudden change

3.Rekey every time the counter reaches its maximum value (If using OTAA: go through OTAA activation procedure again)

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Eavesdropping

Reason for this attack: the block cipher (AES) recreates exactly the same key material every time the counter values repeat.

Specifically:

(plaintext P) XOR (key stream K) = ciphertext C

When given 2 plaintexts P1 and P2 encrypted under the same keystream K, then we have

$$C_1 \oplus C_2 = (P_1 \oplus K) \oplus (P_2 \oplus K)$$
$$= P_1 \oplus P_2 \oplus \underbrace{(K \oplus K)}_{\text{cancels out}}$$
$$= P_1 \oplus P_2.$$

Eavesdropping Cont'd

Step 1: Guess a part of the content in P1

Step 2: Derive the part of P2 at the corresponding position

If all the plaintexts are readable, the guess is possibly correct

Validation and 'enhance' learning: the more resets, the more likely it is to recover the message

Eavesdropping Mitigation

It's the counter value's mistake again.

1. Replace the counter value by a nonce, which is generated from a cryptographically-secure pseudo random number generator

2. Rekey on reset (easier to achieve by OTAA as well)

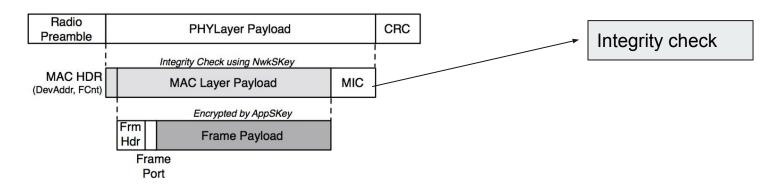
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Bit-flipping Attack

There should have been *integrity checking* to prevent this from happening. But why?

Reason: Encryption and integrity check do NOT happen at the same scope.

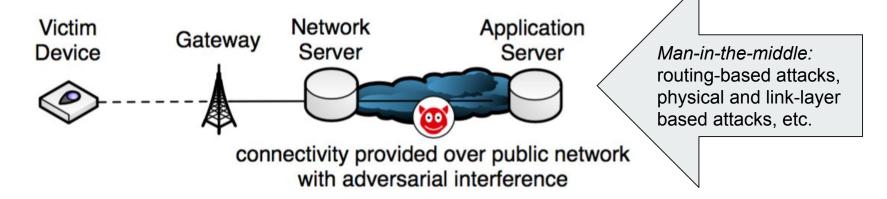
Recall this figure:



Bit-flipping Attack Cont'd

In between the infrastructure operator's *network server* and the IoT solution provider's *application server*, the content cannot be checked for integrity and authenticity

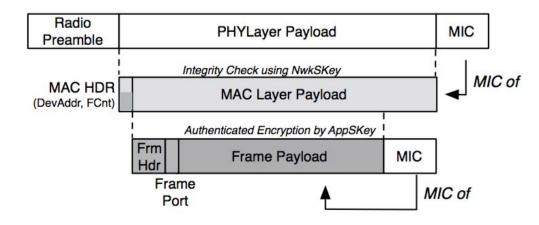
i.e. The integrity check is terminated too early and there is no integrity check throughout transmission



Bit-flipping Attack Mitigation

1. Run the integrity check value at the *application server* instead of the network server

2.Repurpose protocol fields: replace the CRC by a MIC



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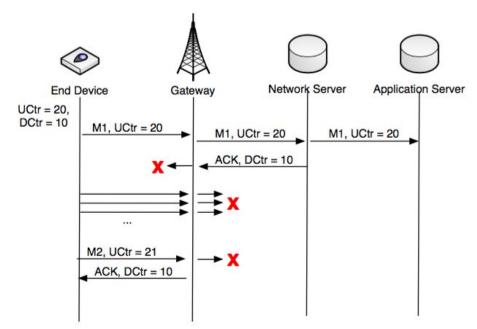
ACK Spoofing

To maximize battery life and reduce the time the radio needs to be powered up, ACK messages from network server do NOT specify which message it is confirming:

PHYSICAL PAYLOAD FORMAT OF AN ACK MESSAGE.

MHDR	DevAddr	FCtrl	FCnt	MIC		
60	88889999	20	0B00	BAE1557A		
	seque	sequential number of all downstream				
	messa	messages				

ACK Spoofing Cont'd



Assume that the gateway is malicious, selectively suppressing certain frames from transmission.

First it blocks the ACK from network server, making end device believe that M1 fails to transfer.

Then when end device send M2, this malicious gateway responds using ACK it blocked before to fool end device, making the end device believe M2 is successfully delivered.

ACK Spoofing Mitigation

ACK spoofing happens because the ACK doesn't specify which message it actually confirms.

Recall in *Bit-flipping Attack Mitigation* part that we use MIC to guarantee integrity throughout the whole transmission.

Apply MIC to both the connection to the network server and the application server:

- then it is possible to add a cryptographic checksum with the returned ACK, including the entire packet as sent by the field device
- IoT device can confirm that ACK belongs to this message and the value of the message remains unchanged during the transmission

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LoRa Class B Attacks

LoRa Class B network: field devices periodically wake up to wait for any incoming messages, the durations of which are specified by the *beacons broadcast* by the gateway

Class B network balances power consumption and the possibility to periodically relay downlink instructions.

Problem: *beacons* are not encrypted, nor protected against malicious modification.

- Spoofing the location of a LoRa gateway (location information can be modified)
- Exhausting the battery

LoRa Class B Attacks Mitigation

Background: beacon frames are not protected, and are especially lacking in integrity check value.

Solution: also change the PHY CRC to a MIC \rightarrow authenticating the beacon frames

